

Department of Computer Science CSCI 2824: Discrete Structures Chris Ketelsen

> Lecture 6: Nested Quantifiers



#### Last Time:

- Introduced predicates and propositional functions
- Started on universal and existential quantifiers

#### **Universal Quantifier:**

•  $\forall x P(x)$ : "For all x in my domain P(x) is true "

#### **Existential Quantifier:**

•  $\exists x \ P(x)$ : "There exists an x in my domain s.t. P(x) is true"



**Warm-Up Problems**: Let the domain for x be the set of all Natural Numbers  $\mathbb{N} = \{0, 1, 2, \dots\}$ 

Numbers, 
$$\mathbb{N} = \{0, 1, 2, ...\}$$

**Example**: Determine the truth value of  $\forall n \ (3n \le 4n)$ 

**Warm-Up Problems**: Let the domain for x be the set of all Natural Numbers,  $\mathbb{N} = \{0, 1, 2, ...\}$ 

**Example**: Determine the truth value of  $\forall n \ (3n \le 4n)$ 

This is true. Q: What if the domain was the set of all integers?

**Example**: Determine the truth value of  $\exists x (x^2 = x)$ 

**Warm-Up Problems**: Let the domain for x be the set of all Natural Numbers,  $\mathbb{N} = \{0, 1, 2, ...\}$ 

**Example**: Determine the truth value of  $\forall n \ (3n \le 4n)$ 

This is true. Q: What if the domain was the set of all integers?

**Example**: Determine the truth value of  $\exists x (x^2 = x)$ 

This is true. We just need to find one x in the domain that works, and in this case there are two: x = 0 and x = 1.

Last time we showed the following equivalences

### DeMorgan's Laws for Quantifiers:

- $\neg \forall x P(x) \equiv \exists x \neg P(x)$
- $\neg \exists x P(x) \equiv \forall x \neg P(x)$

#### Distribution Laws for Quantifiers:

- $\forall x (P(x) \land Q(x)) \equiv \forall x P(x) \land \forall x Q(x)$
- $\exists x (P(x) \lor Q(x)) \equiv \exists x P(x) \lor \exists x Q(x)$

**Note**: Distribution of ∀ over ∨ and ∃ over ∧ didn't work

### A Computer Sciency Way of Viewing Quantifiers

Think of quantified statements as loops that do logic checks

Example:  $\forall x \ P(x)$ 

```
In [ ]: for x in domain:
    if P(x) == False:
        return False
    return True
```

- If we find an x in domain where P(x) is False, return False
- If we make it through loop then return True

### A Computer Sciency Way of Viewing Quantifiers

Think of quantified statements as loops that do logic checks

Example:  $\exists x \ P(x)$ 

```
In [ ]: for x in domain:
    if P(x) == True:
        return True
    return False
```

- If we find an x in domain where P(x) is True, return True
- If we make it through loop without finding one, return False

#### **Nested Quantifiers**

Interesting things happen when we include multiple quantifiers

**Example**: What does this say:  $\forall x \exists y (x + y = 0)$ ?

$$Q(x) = \exists y (x+y=0)$$

$$\exists x \forall y (x + y = 0)$$

#### **Nested Quantifiers**

Interesting things happen when we include multiple quantifiers

**Example**: What does this say:  $\forall x \exists y (x + y = 0)$ ?

It really helps to read these outloud: "For all x, there exists a y, such that the sum of x and y is zero"

What do you think? Is this true or false?

#### **Nested Quantifiers**

Interesting things happen when we include multiple quantifiers

**Example**: What does this say:  $\forall x \exists y (x + y = 0)$ ?

It really helps to read these outloud: "For all x, there exists a y, such that the sum of x and y is zero"

What do you think? Is this true or false?

This is totally **true**. It's the expression of the fact that all numbers have an **additive inverse**.

### **Nested Quantifiers as Loops**

**Example**:  $\forall x \exists y P(x, y)$  ?

- If we make it through y-loop without finding a True, return
   False
- If we make it through entire x-loop then return True



### **Nested Quantifiers as Loops**

**Example**:  $\forall x \exists y (x + y = 0)$ ?

```
In [7]: def check_additive_inverse(domain):
    for x in domain:
        exists_y = False
        for y in domain:
            if x + y == 0:
                exists_y = True
        if exists_y == False:
            return False
    return True

domain = [-3, -2, -1, 0, 1, 2, 3]
    check_additive_inverse(domain)
```

### **Nested Quantifiers as Loops**

**Example**:  $\forall x \exists y (x + y = 0)$ ?

```
In [8]: def check_additive_inverse(domain):
    for x in domain:
        exists_y = False
        for y in domain:
            if x + y == 0:
                 exists_y = True
        if exists_y == False:
            return False
    return True

domain = [-2, -1, 0, 1, 2, 3]
    check_additive_inverse(domain)
```

Out[8]: False

**Example**: How could we express the law of **commutation of** addition (that is, that x + y = y + x)?

**Example**: How could we express the law of **commutation of** addition (that is, that x + y = y + x)?

How about  $\forall x \ \forall y \ (x + y = y + x)$ 

### **Nested Quantifiers as Loops**

**Example**:  $\forall x \ \forall y \ P(x, y)$  ?

- If we ever find an (x, y)-pair that makes P(x, y) False, return False
- If we make it through both loops, return True

### **Nested Quantifiers as Loops**

**EFY**: Cook up an example of a statement of the form  $\exists x \forall y \ P(x, y)$  that is True, and write a Python function with nested for-loops that checks it!

**EFY**: Cook up an example of a statement of the form  $\exists x \exists y \ Q(x, y)$  that is True, and write a Python function with nested for-loops that checks it!

**Example**: How could we express the law of **commutation of** addition (that is, that x + y = y + x)?

How about  $\forall x \ \forall y \ (x + y = y + x)$ 

**Question**: What happens if we change the order of  $\forall x$  and  $\forall y$ ?

**Example**: How could we express the law of **commutation of** addition (that is, that x + y = y + x)?

How about  $\forall x \ \forall y \ (x + y = y + x)$ 

**Question**: What happens if we change the order of  $\forall x$  and  $\forall y$ ?

So we'd have  $\forall y \ \forall x \ (x + y = y + x)$ 

**Answer**: Not much! Still looping over all pairs of x's and y's

Let's go back to the previous example:

**Example**:  $\forall x \exists y (x + y = 0)$ 

Question: What happens if we change the order here?

Let's go back to the previous example:

**Example**:  $\forall x \ \exists y \ (x + y = 0)$ 

Question: What happens if we change the order here?

**Answer**: A lot! The new expression  $\exists y \ \forall x \ (x + y = 0)$  says

• "There exists some number y such that for every x out there, x+y=0"

Can you think of such a number?

Let's go back to the previous example:

**Example**:  $\forall x \ \exists y \ (x + y = 0)$ 

Question: What happens if we change the order here?

**Answer**: A lot! The new expression  $\exists y \ \forall x \ (x + y = 0)$  says

• "There exists some number y such that for every x out there, x+y=0"

Can you think of such a number?

Me neither! In fact, after switching the order of the quantifiers the proposition becomes false.

### Rules for Switching Quantifiers:

- OK to switch  $\forall x$  and  $\forall y$
- OK to switch  $\exists x$  and  $\exists y$  (**EFY**: Check that this is true!)
- **NOT** OK to switch  $\forall x$  and  $\exists y$

OK, let's do a bunch more examples

**Example:** Now we'll switch the domain to all real numbers

How can you express the fact that all numbers of have a multiplicative inverse

That is, a number that you can multiply by to get 1?

First of all, is it really true that **all** numbers of have a multiplicative

Inverse?

1: Domain = 
$$\mathbb{Z} - \xi o \xi$$
  $\forall x \neq y (x y = 1)$ 

2: Domain =  $\mathbb{Z}$   $\forall x \neq y (x y = 1) \lor (x = 0)$ 

\*  $\forall x \neq y (x y = 1) \lor (x = 0)$ 

\*  $\forall x \neq y (x y = 1) \lor (x = 0)$ 

Example: Now we'll switch the domain to all real numbers

How can you express the fact that all numbers of have a multiplicative inverse

That is, a number that you can multiply by to get 1?

First of all, is it really true that **all** numbers of have a multiplicative inverse?

Nope! But all nonzero numbers do

So how could we say this with quantifiers?

Let's say it in logic-y English

"For all x's that aren't zero, there exists a y such that xy = 1"

Let's say it in logic-y English

"For all x's that aren't zero, there exists a y such that xy = 1"

Note that x not being zero is a **condition** that has to happen before we consider looking for an inverse. Let's rephrase:

"For all x, if  $x \neq 0$  then there exists a y such that xy = 1"

How about

$$\forall x ((x \neq 0) \rightarrow \exists y (xy = 1))$$

**Example**: How could you express that there are an infinite number of natural numbers?

$$\forall x \exists y (y = x+1) \forall x \exists y (y > x)$$

**Example**: How could you express that there are an infinite number of natural numbers?

If domains for x and y are the set of natural numbers, we could say

$$\forall x \; \exists y \; (y > x)$$

This just says that every natural number has a number that is larger

**EFY**: How could you express that if you multiply two negative numbers together you get a positive number?

**EFY**: How could you express that the real numbers have a **multiplicative identity**. That is, that there's a number out there that when you multiply something by it, you get the same thing back. (**Note**: this is literally saying that the number 1 is a thing)

OK, lets practice some non-mathy translations

**Example:** Translate the statement "You can fool some of the people all of the time"

We need to define a propositional function that says a person can be fooled at a particular time

Let F(p, t) represent "you can fool person p at time t"

Then we have  $\exists p \ \forall t \ F(p, t)$ 

**Example:** Translate the statement "You can fool all of the people some of the time"

To me, this one is actually kinda ambiguous.

Does it mean "There is a time when you can fool all of the people"?

In which case we would have  $\exists t \ \forall p \ F(p, t)$ 

Or does it mean "Each person has a time that they could be fooled"?

In which case we would have  $\forall p \; \exists t \; F(p, t)$ 

Rule of Thumb: Logic and mathematics are precise but language ISN'T so you have to be cautions



**Example:** Translate the statement "You can't fool all of the people all of the time"

This works out to be  $\neg(\forall p \ \forall t \ F(p, t))$ 

What would happen if we pushed the negation through?

$$\neg \forall p \ \forall t \ F(p,t) \equiv \exists p \ \neg \forall t \ F(p,t) \equiv \exists p \ \exists t \ \neg F(p,t)$$

which translates to the equivalent (but more awkward) statement

"There is some person at some time that can't be fooled"

Quantifications with more than two quantifiers are also common

**Example**: Let Q(x, y, z) mean "x + y = z". What are the truth values of

- $\forall x \ \forall y \ \exists z \ Q(x, y, z)$
- $\exists z \ \forall x \ \exists y \ Q(x, y, z)$

**Example**: One more! Translate the following using quantifiers:

Babies are illogical. Nobody is despised who can manage a crocodile. Illogical people are despised. Therefore, babies cannot manage crocodiles

Let B(x) mean "x is a baby", L(x) mean "x is logical", C(x) mean "x can handle a crocodile", and D(x) mean "x is despised"

### **End of Representational Logic**

- We now know how to represent standard propositions
- We know how to represent propositions with quantifiers
- We know how to prove and derive logical equivalences

### **Next Time We Start Learning to Argue**

- Rules of inference
- Valid and sound arguments
- Proof types and strategies

**EFY**: Cook up an example of the form  $\exists x \ \forall y \ P(x, y)$  that is True, and write a Python function with nested for-loops that checks it!

**Solution**: How about  $\exists x \ \forall y \ xy = 0$  (essentially, 0 exists)

**EFY**: Cook up an example of the form  $\exists x \exists y \ Q(x, y)$  that is True, and write a Python function with nested for-loops that checks it!

**Solution**: How about  $\exists x \exists y \ x^2 + y^2 = 25$ 

```
In [15]: def check_sum_of_squares(domain):
    for x in domain:
        for y in domain:
            if x**2 + y**2 == 25:
                return True

    return False

domain = [ 0, 1, 2, 3, 4, 5]
    check_sum_of_squares(domain)
```

**EFY**: Is it OK to switch the order of  $\exists x \ \exists y$ ?

**Solution**: Totally. Consider the example "There exists an integer x and an integer y such that  $x^2 + y^2 = 25$ ".

This is true because we can let x = 3 and y = 4

**Question**: What changes if we write it as "There exists an integer y and an integer x such that  $x^2 + y^2 = 25$ "?

**Answer**: Literally nothing

**EFY**: How could you express that if you multiply two negative numbers together you get a positive number?

**Solution**: We want to say that if we take any pair of numbers, if those numbers are negative their product is positive.

How about

$$\forall x \ \forall y (((x < 0) \land (y < 0)) \rightarrow (xy > 0))$$

**EFY**: How could you express that the real numbers have a **multiplicative identity**. That is, that there's a number out there that when you multiply something by it, you get the same thing back. (**Note**: this is literally saying that the number 1 is a thing)

Solution: We want to say

"There exists a number such that for any x when you multiply that number by x the result is x"

How about

$$\exists y \ \forall x \ (xy = x)$$