

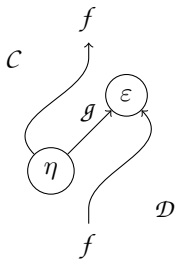
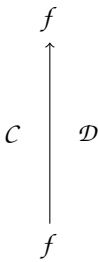
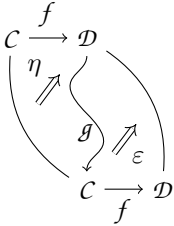
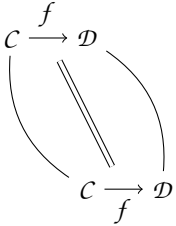
# Adjunctions

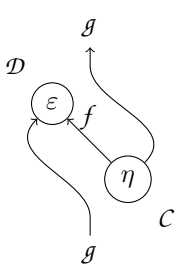
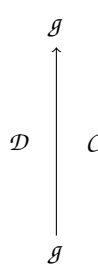
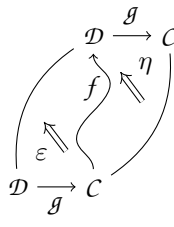
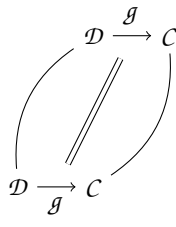
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**Definition** (Adjunction for a 2-Category  $\mathbf{C}$ ). A tuple  $\langle C, \mathcal{D}, f, g, \eta, \varepsilon, \mathfrak{f}, \mathfrak{g} \rangle$  whose components have the following types:

$C$  is an object of  $\mathbf{C}$   
 $\mathcal{D}$  is an object of  $\mathbf{C}$   
 $f$  is a morphism of  $\mathbf{C}$  from  $C$  to  $\mathcal{D}$   
 $g$  is a morphism of  $\mathbf{C}$  from  $\mathcal{D}$  to  $C$   
 $\eta$  is a 2-cell of  $\mathbf{C}$  from  $C$  to  $f;g$   
 $\varepsilon$  is a 2-cell of  $\mathbf{C}$  from  $g;f$  to  $\mathcal{D}$

$\mathfrak{f}$  is a proof that  equals . In other words,  equals .

$\mathfrak{g}$  is a proof that  equals . In other words,  equals .

**Definition** (Left/Right Adjoint).  $f$  above is called the left adjoint, and  $g$  above is called the right adjoint. A morphism of a 2-category is a left/right adjoint if it is the left/right adjoint of some adjunction.

**Exercise 1.** Prove that there is a bijection between adjunctions in  $\mathbf{Cat}$  and adjunctions via transpositions.

**Example.** Consider  $\mathbf{Prost}$ . Suppose we had a pair of preordered sets  $\langle C, \leq \rangle$  and  $\langle D, \leq \rangle$ , and we want to make an adjunction out of some relation-preserving functions  $f : C \rightarrow D$  and  $g : D \rightarrow C$ . Then  $\eta$  exists if and only if  $\forall c : C. c \leq g(f(c))$ , and  $\beta$  exists if and only if  $\forall d : D. f(g(d)) \leq g$ . If  $\eta$  and  $\beta$  exist, then  $\mathfrak{f}$  and  $\mathfrak{g}$  are trivial since  $\mathbf{Prost}$  is a *locally thin* 2-category. Such a situation is called a monotone Galois connection.