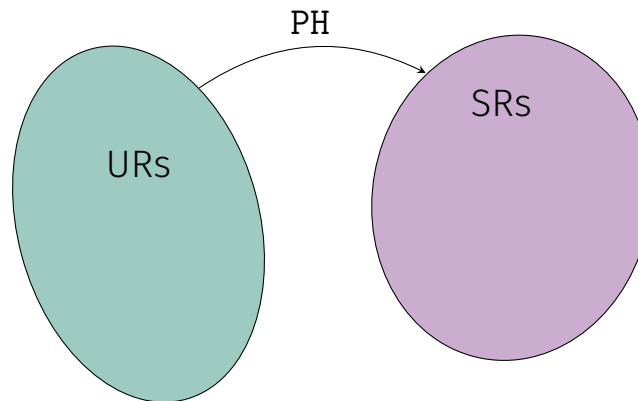


A subregular approach to the problem of learning underlying representations

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December 9, 2019 · Tel Aviv University

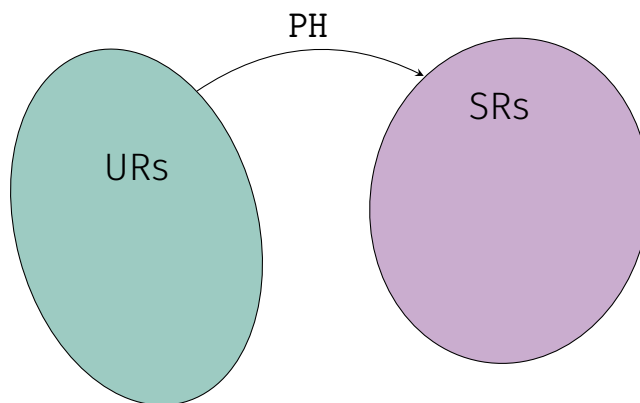
What is the nature of phonology?



What is the nature of:

- constraints on SRs? URs? (Halle, 1959; Prince and Smolensky, 1993; Gorman, 2013)
- maps from URs to SRs? (Chomsky and Halle, 1968; Johnson, 1972)
- relation between SRs and URs?
(Hyman, 1970; Kiparsky, 1973; Kenstowicz and Kisseberth, 1977)

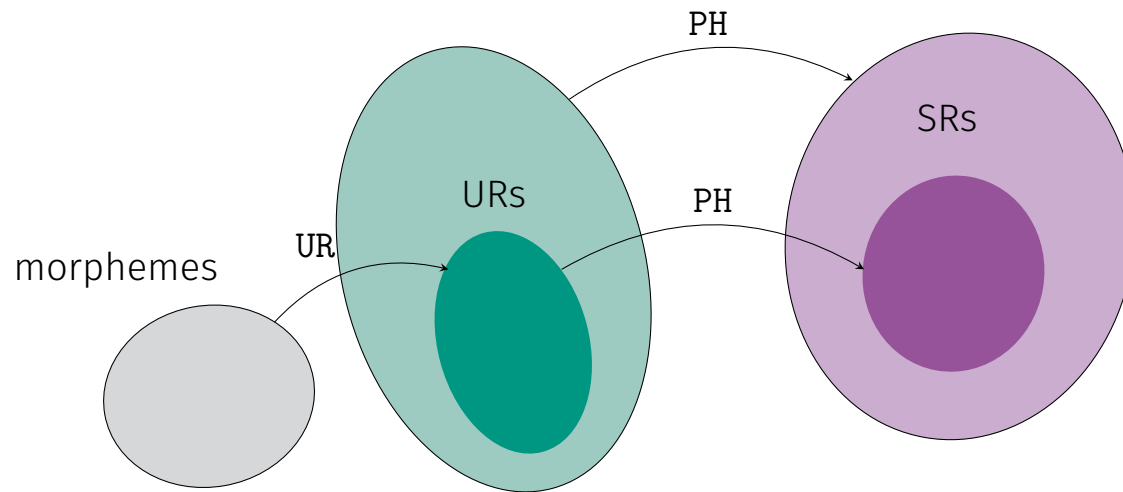
What is the nature of phonology?



What is the **computational** nature of **(learning)**

- **constraints on SRs?** URs? (Heinz, 2009, 2010)
- **maps from URs to SRs?** (Jardine et al., 2014; Chandlee and Heinz, 2018)
- relation between SRs and URs?

Learning URs and a grammar



- Computational restrictions on maps from URs to SRs provide avenue for **learning URs and a grammar**
- This includes **restrictions on relation between SRs and URs**

Learning URs and a grammar

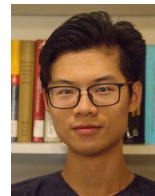
- **Learning problem:** the **simultaneous inference** of URs and a grammar from SRs in a morphological paradigm
(Tesar, 2014; Cotterell et al., 2015; Rasin et al., 2018)
- **Today:** the **subsequential** functions provide a structure that can solve this problem
(Mohri, 1997; Heinz and Lai, 2013; Jardine et al., 2014)
- More specifically, we'll look at **input strictly local (ISL)** functions
(Chandlee and Heinz, 2018)

Learning URs and a grammar

- This poses further restrictions on the relationship between SRs and URs
- This is joint work with students at Rutgers



Wenyue Hua



Huteng Dai

- This is very much work in progress!

The learning problem

The learning problem

English plural:

CAT-PL	[kæt ^s]
CUFF-PL	[kʌf ^s]
DEATH-PL	[dεθ ^s]
GIRL-PL	[gərl ^z]
CHAIR-PL	[tʃer ^z]
BOY-PL	[bɔɪ ^z]
...	...

Analysis:

- A **map** from morphemes to URs

CAT	→	/kæt/
PL	→	/z/
...		...

- A **map** from URs to SRs

/z/ → [s] / [-voi] _____

The learning problem

- M : finite set of morphemes
- Σ : finite set of segments
- Learning targets:
 - lexicon function $UR : M^* \rightarrow \Sigma^*$

{CAT, DOG, ..., PL}

{a, b, β, ..., z}

- phonology function $PH : \Sigma^* \rightarrow \Sigma^*$

$UR(CAT) = kæt$

$UR(PL) = z$

$UR(CAT-PL) = kætz$

...

$PH(kæt) = kæt$

$PH(dɔgz) = dɔgz$

$PH(kætz) = kæts$

$PH(bnɪkz) = bnɪks$

...

The learning problem

- Learning data is a finite sample of $\text{PH} \circ \text{UR}$

$w \in M^*$	$\text{PH}(\text{UR}(w))$
CAT-PL	kæts
CUFF-PL	kʌfs
DEATH-PL	dεθs
GIRL-PL	gərlz
CHAIR-PL	tʃerz
BOY-PL	bɔɪz

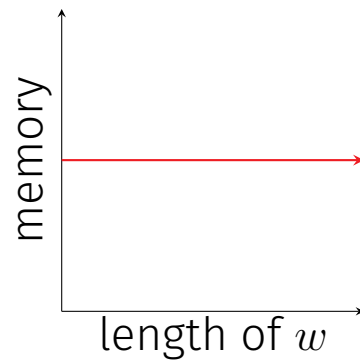
The learning problem

- **Problem:** identify UR and PH from a finite sample of $\text{PH} \circ \text{UR}$
- **Questions:** What is the nature of ...
 - UR
 - PH
 - $\text{PH} \circ \text{UR}$
 - the data sample...such that learning is possible?

Subsequentiality and phonology

Subsequentiality and phonology

- Johnson (1972); Kaplan and Kay (1994): phonological maps are **regular**

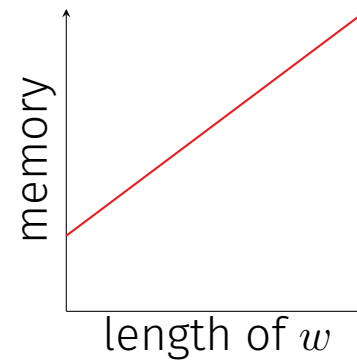


regular

E.g., directional harmony

/k*a*k*i*-kæ/ \mapsto [k*a*k*i*-k*a*]

/k*i*k*i*-kæ/ \mapsto [k*i*k*i*-k*æ*]



non-regular

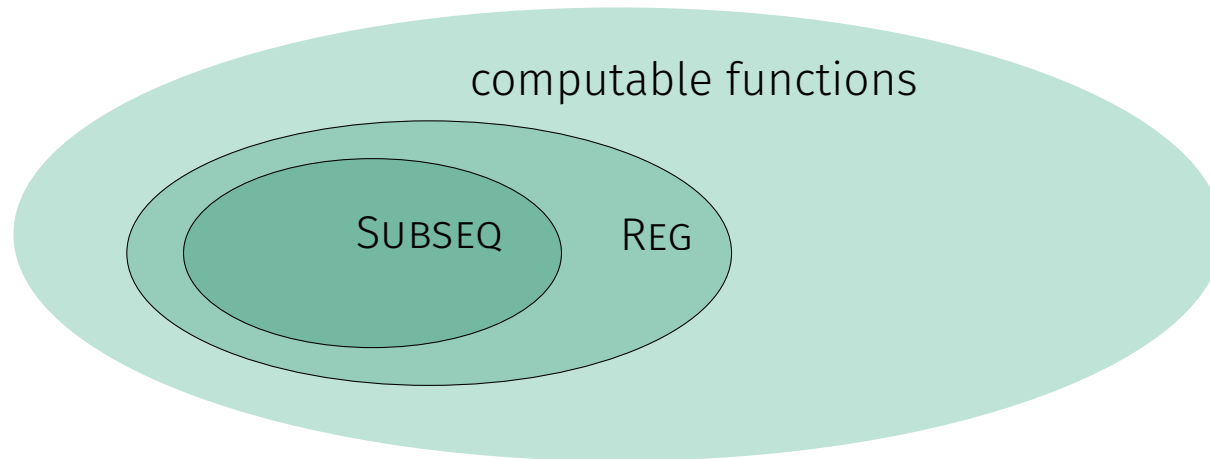
E.g., “majority rules” (Baković, 2000)

/k*a*k*a*-kæ/ \mapsto [k*a*k*a*-k*a*]

/k*æ*k*a*-kæ/ \mapsto [k*æ*k*æ*-k*æ*]

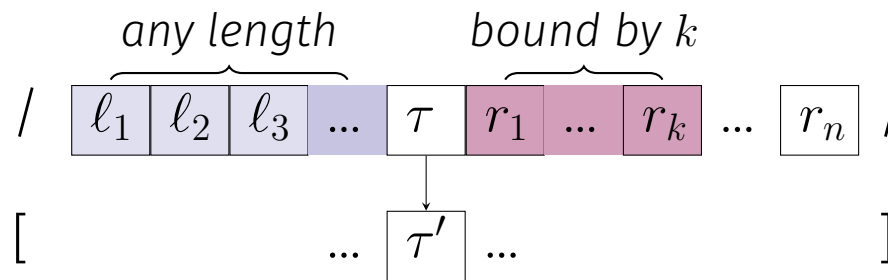
Subsequentiality and phonology

- Mohri (1997); Heinz and Lai (2013): Phonological maps are **subsequential**;
 - they are regular, **and**
 - they are **deterministic**



Subsequentiality and phonology

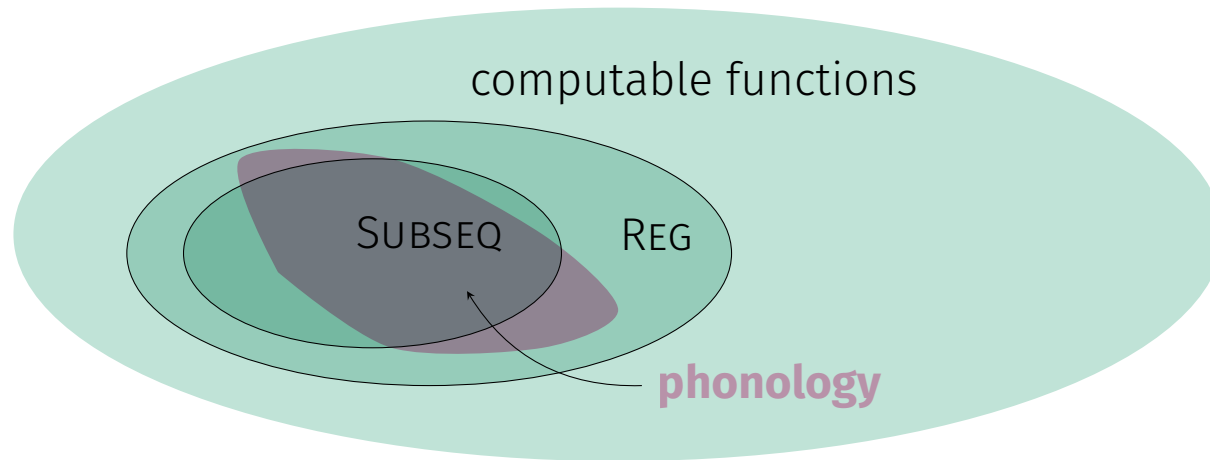
- **Subsequential:** output can be determined deterministically in one direction



- (Deterministic \neq no optionality; Heinz in progress)

Subsequentiality and phonology

- The subsequentiality of phonology is empirically well-supported (Chandlee and Heinz, 2012; Heinz and Lai, 2013; Payne, 2017; Luo, 2017; Chandlee and Heinz, 2018)



- Though c.f. Jardine (2016); McCollum et al. (2017)

Subsequentiality and phonology

- The **longest common prefix (lcp)** is the longest initial sequence shared by a set of strings

$$\text{lcp}(\{aab, aaba, aac\}) =$$

Subsequentiality and phonology

- The **longest common prefix (lcp)** is the longest initial sequence shared by a set of strings

$$\text{lcp}(\{\underline{a}ab, \underline{a}aba, \underline{a}ac\}) = aa$$

Subsequentiality and phonology

- The **longest common prefix (lcp)** is the longest initial sequence shared by a set of strings

$$\text{lcp}(\{\underline{a}ab, \underline{a}aba, \underline{a}ac\}) = aa$$

$$\text{lcp}(\{bac, abc\}) =$$

Subsequentiality and phonology

- The **longest common prefix (lcp)** is the longest initial sequence shared by a set of strings

$$\text{lcp}(\{\underline{aab}, \underline{aaba}, \underline{aac}\}) = aa$$

$$\text{lcp}(\{bac, abc\}) = \lambda$$

Subsequentiality and phonology

- Take a function f

$$\left. \begin{array}{l} f(tat) = tat \\ f(tatta) = tatta \\ f(tadta) = tadda \\ f(dtd) = ddd \\ f(tadtta) = taddtta \\ \dots \end{array} \right\} t \rightarrow d / d _ \text{(simul.)}$$

Subsequentiality and phonology

- For f we define

$$f^p(w) \stackrel{\text{def}}{=} \text{1cp}(f(w\Sigma^*))$$

Subsequentiality and phonology

- For f we define

$$f^p(w) \stackrel{\text{def}}{=} \text{1cp}(f(w\Sigma^*))$$

- Let $\Sigma = \{a, d, t\}$

$$f^p(tad) = \text{1cp}(\{ f(tad), \\ f(tada), \\ f(tadd), \\ f(tadt), \\ f(tadaa), \\ \dots \})$$

Subsequentiality and phonology

- For f we define

$$f^p(w) \stackrel{\text{def}}{=} \text{1cp}(f(w\Sigma^*))$$

- Let $\Sigma = \{a, d, t\}$

$$f^p(tad) = \text{1cp}(\{ \begin{array}{l} f(tad) = tad, \\ f(tada) = tada, \\ f(tadd) = tadd, \\ f(tadt) = tadd, \\ f(tadaa) = tadaa, \\ \dots \end{array} \})$$

Subsequentiality and phonology

- For f we define

$$f^p(w) \stackrel{\text{def}}{=} \text{1cp}(f(w\Sigma^*))$$

- Let $\Sigma = \{a, d, t\}$

$$f^p(tad) = \text{1cp}(\{ \begin{array}{l} f(tad) = \underline{tad}, \\ f(tada) = \underline{tada}, \\ f(tadd) = \underline{tadd}, \\ f(tadt) = \underline{tadd}, \\ f(tadaa) = \underline{tadaa}, \\ \dots \end{array} \})$$

Subsequentiality and phonology

- For f we define

$$f^p(w) \stackrel{\text{def}}{=} \text{1cp}(f(w\Sigma^*))$$

- Let $\Sigma = \{a, d, t\}$

$$f^p(tad) = tad$$

Subsequentiality and phonology

- For f we define

$$f^p(w) \stackrel{\text{def}}{=} \text{1cp}(f(w\Sigma^*))$$

- Let $\Sigma = \{a, d, t\}$

$$\begin{aligned} f^p(tad) &= tad \\ f^p(tadt) &= \text{1cp}(\{ f(tadt) = \underline{tadd}, \\ &\quad f(tadta) = \underline{tadda}, \\ &\quad f(tadtd) = \underline{taddd}, \\ &\quad f(tadtt) = \underline{taddt}, \\ &\quad f(tadtaa) = \underline{taddaa}, \\ &\quad \dots \quad \quad \quad \}) \end{aligned}$$

Subsequentiality and phonology

- For f we define

$$f^p(w) \stackrel{\text{def}}{=} \text{1cp}(f(w\Sigma^*))$$

- Let $\Sigma = \{a, d, t\}$

$$\begin{aligned} f^p(tad) &= tad \\ f^p(tadt) &= tadd \end{aligned}$$

Subsequentiality and phonology

- $f^p(w)$ is the **contribution** of w to any $f(wv)$

$$f(wv) = \overbrace{\boxed{a} \boxed{b} \boxed{a} \dots \boxed{a} \boxed{b} \boxed{a}}^{f^p(w)} \dots \boxed{b}$$

- f^p grows proportionally iff f subsequential...

$$\begin{aligned} m^p(\text{kækak}) &= \text{1cp}(\{\underline{\text{k}}\text{ækækæ}, \underline{\text{k}}\text{akaka}, \dots\}) \\ &= \text{k} \end{aligned}$$

Subsequentiality and phonology

- For $w \in \Sigma^*$, the **environment function** is

$$f_w(v) \stackrel{\text{def}}{=} f^p(w)^{-1} f(wv)$$

$$f(wv) = \begin{array}{|c|c|c|c|c|c|c|c|c|} \hline & \overbrace{a \quad b \quad a \quad \dots \quad a}^{f^p(w)} & \overbrace{b \quad a \quad \dots \quad b}^{f_w(v)} & \\ \hline \end{array}$$

Subsequentiality and phonology

- For $w \in \Sigma^*$, the **environment function** is

$$f_w(v) \stackrel{\text{def}}{=} f^p(w)^{-1} f(wv)$$

$$f(wv) = \overbrace{\begin{array}{|c|c|c|c|} \hline a & b & a & \dots & a \\ \hline \end{array}}^{f^p(w)} \overbrace{\begin{array}{|c|c|c|c|} \hline b & a & \dots & b \\ \hline \end{array}}^{f_w(v)}$$

- Ex.

$$\begin{aligned} f_{tad}(ta) &= f^p(tad)^{-1} f(tadta) \\ &= (tad)^{-1} \underline{tad}da \\ &= da \end{aligned}$$

Subsequentiality and phonology

- For $w \in \Sigma^*$, the **environment function** is

$$f_w(v) \stackrel{\text{def}}{=} f^p(w)^{-1} f(wv)$$

$$f(wv) = \overbrace{a \ b \ a \ \dots \ a}^{f^p(w)} \overbrace{b \ a \ \dots \ b}^{f_w(v)}$$

- Ex.

$$\begin{aligned} f_{tad}(ta) &= f^p(tad)^{-1} f(tadta) \\ &= (tad)^{-1} tadda \\ &= da \end{aligned}$$

$$\begin{aligned} f_{tat}(ta) &= f^p(tat)^{-1} f(tatta) \\ &= (tat)^{-1} \underline{tatta} \\ &= ta \end{aligned}$$

Subsequentiality and phonology

- f is subsequential **iff it has finite environment functions**

w	$f_w(ta)$	w	$f_w(ta)$
a	ta	d	da
t	ta	ad	da
aa	ta	dd	da
at	ta	td	da
dt	ta	aad	da
tt	ta	add	da
aaa	ta	\dots	\dots
\dots	\dots		

Subsequentiality and phonology

- f is subsequential **iff it has finite environment functions**

w	$f_w(ta)$	w	$f_w(ta)$
a	ta	d	da
t	ta	ad	da
aa	ta	dd	da
\dots	\dots	\dots	\dots

$$f_a = f_t = \dots = f_{aaa} = \dots = f_{tatat} = \dots = f_{a/t}$$

$$f_d = f_{ad} = \dots = f_{aad} = \dots = f_{tatad} = \dots = f_d$$

Subsequentiality and phonology

- Environment functions correspond to states in a **deterministic finite-state transducer** (Mohri, 1997)
- There are procedures for determining environment functions from positive data (Oncina et al., 1993; Chandlee et al., 2014; Jardine et al., 2014)
- If f 's environment functions represent $k - 1$ suffixes, f is **input strictly k -local (ISL $_k$)** (Chandlee, 2014; Chandlee and Heinz, 2018)

$$f \rightarrow \{f_{a/t}, f_d\}$$

Towards a solution

Towards a solution

- M : set of morphemes; Σ : finite set of segments
- lexicon function $UR : M^* \rightarrow \Sigma^*$
phonology function $PH : \Sigma^* \rightarrow \Sigma^*$
- **Problem:** identify UR and PH from a finite sample of $PH \circ UR$

Towards a solution

- Example:

UR	PH
$r_1 \mapsto \text{tat} \quad s_1 \mapsto \text{ta}$	
$r_2 \mapsto \text{tad} \quad s_2 \mapsto \text{da}$	$t \rightarrow d / d \text{ ___}$
$r_3 \mapsto \text{a} \quad s_3 \mapsto \text{a}$	

Sample of $\text{PH} \circ \text{UR}$					
w	$\text{PH}(\text{UR}(w))$	w	$\text{PH}(\text{UR}(w))$	w	$\text{PH}(\text{UR}(w))$
$r_1 s_1$	tatta	$r_2 s_1$	tadda	$r_3 s_1$	ata
$r_1 s_2$	tatda	$r_2 s_2$	tadda	$r_3 s_2$	ada
$r_1 s_3$	tata	$r_2 s_3$	tada	$r_3 s_3$	aa

Towards a solution

- M : set of morphemes; Σ : finite set of segments
- lexicon function $UR : M^* \rightarrow \Sigma^*$
phonology function $PH : \Sigma^* \rightarrow \Sigma^*$
- **Problem:** identify UR and PH from a finite sample of $PH \circ UR$
- What is the nature of ...
 - UR
 - PH
 - $PH \circ UR$...?

Towards a solution

- M : set of morphemes; Σ : finite set of segments
- lexicon function $UR : M^* \rightarrow \Sigma^*$
phonology function $PH : \Sigma^* \rightarrow \Sigma^*$
- **Problem:** identify UR and PH from a finite sample of $PH \circ UR$
- The nature of ...
 - UR
 - PH
 - $PH \circ UR$...is **subsequential**

Towards a solution

Assumptions:

- UR has **one environment function** ($= \text{UR}$)

$$\text{UR}_w(\text{CAT}) = \text{kæt} \quad \text{for any } w \in M^*;$$

$$\text{UR}_w(\text{PL}) = \text{z} \quad \text{for any } w \in M^*; \text{ etc.}$$

Towards a solution

Assumptions:

- UR has **one environment function** (= UR)

$$\text{UR}_w(\text{CAT}) = \text{kæ}t \text{ for any } w \in M^*;$$

$$\text{UR}_w(\text{PL}) = z \text{ for any } w \in M^*; \text{ etc.}$$

- PH is ISL_2
- That is, its environment functions are of the form PH_σ , $\sigma \in \Sigma$

$$\Sigma = \{t, a, d\} \rightarrow \text{possible env. functions are } \text{PH}_a, \text{PH}_t, \text{PH}_d$$

Towards a solution

Strategy:

- Two hypotheses UR' and PH'
- We modify UR' so it has one environment function
- We make the *opposite* change in PH' to remain consistent with input data

The procedure

The procedure

- Running example $D \subset \text{PH} \circ \text{UR}$

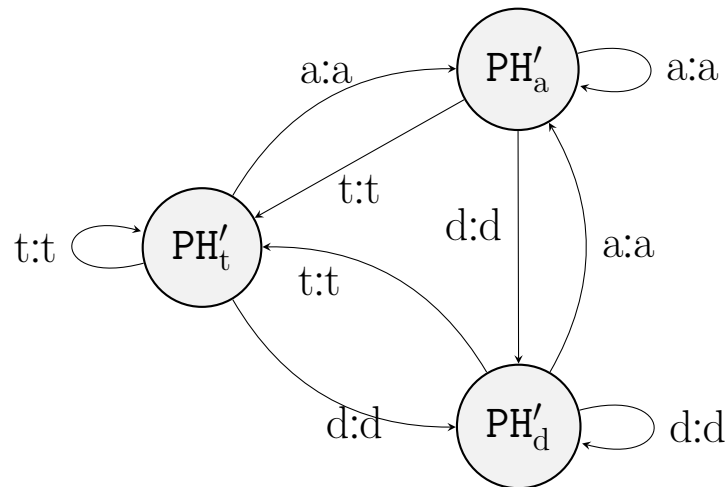
Sample of $\text{PH} \circ \text{UR}$

w	$\text{PH}(\text{UR}(w))$	w	$\text{PH}(\text{UR}(w))$	w	$\text{PH}(\text{UR}(w))$
r_1s_1	tatta	r_2s_1	tadda	r_3s_1	ata
r_1s_2	tatda	r_2s_2	tadda	r_3s_2	ada
r_1s_3	tata	r_2s_3	tada	r_3s_3	aa

The procedure

- Initialize PH' to the identity function

$PH'(tad) = tad$, $PH'(tatta) = tatta$, $PH'(tadta) = tadta$, etc.



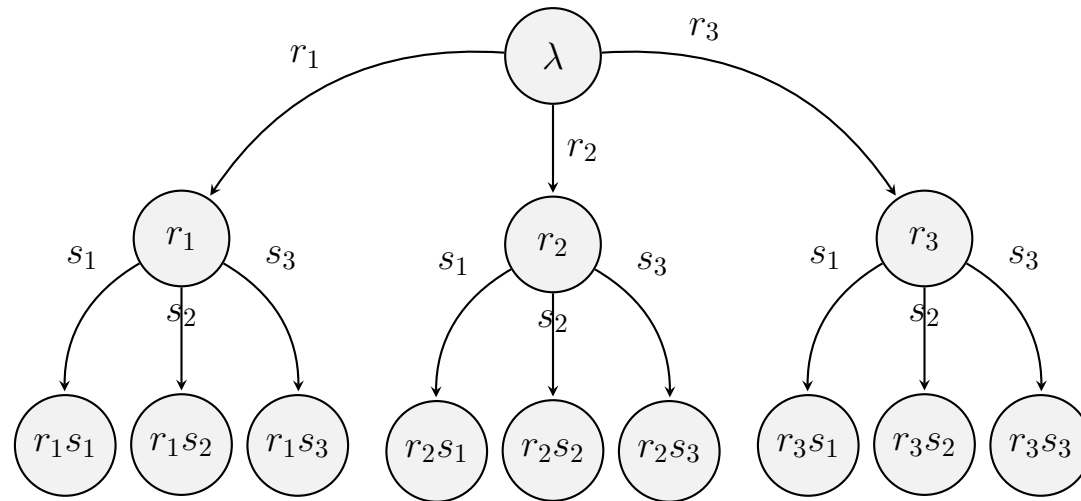
The procedure

- Initialize UR' to a **prefix tree transducer** representing D

r_1s_1	tatta	r_2s_1	tadda	r_3s_1	ata
r_1s_2	tatda	r_2s_2	tadda	r_3s_2	ada
r_1s_3	tata	r_2s_3	tada	r_3s_3	aa

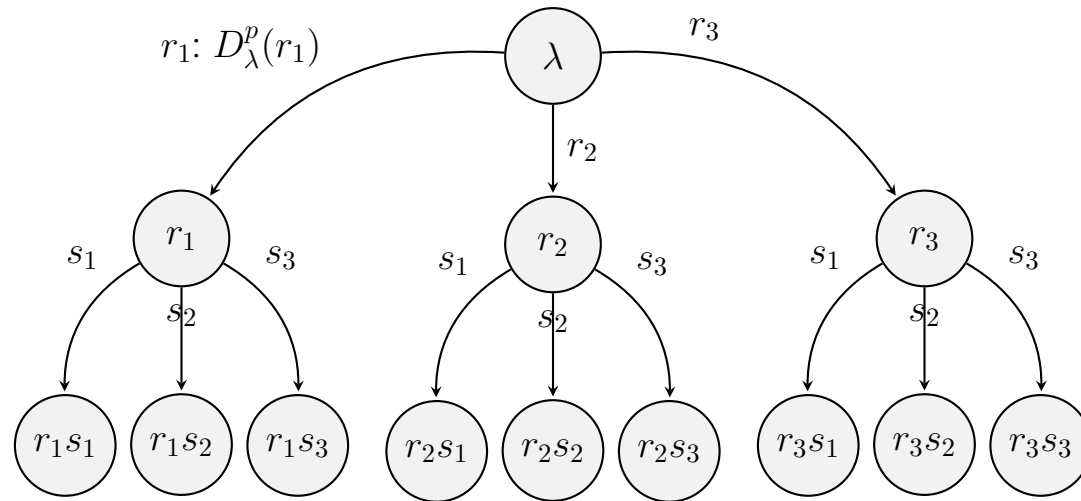
The procedure

<u>r_1s_1</u>	tatta	<u>r_2s_1</u>	tadda	<u>r_3s_1</u>	ata
<u>r_1s_2</u>	tatda	<u>r_2s_2</u>	tadda	<u>r_3s_2</u>	ada
<u>r_1s_3</u>	tata	<u>r_2s_3</u>	tada	<u>r_3s_3</u>	aa



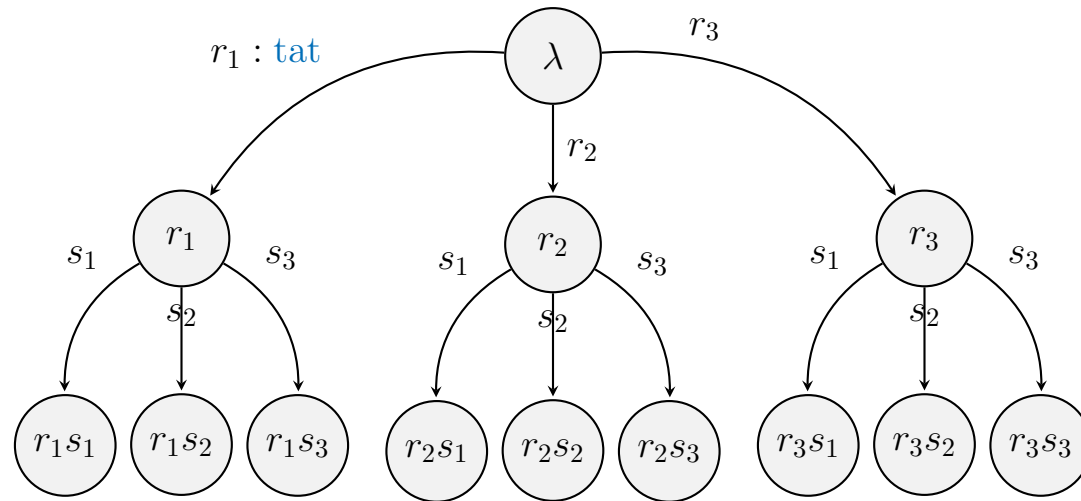
The procedure

$m : D_w^p(m)$	r_1s_1	tatta	r_2s_1	tadda	r_3s_1	ata
	r_1s_2	tatda	r_2s_2	tadda	r_3s_2	ada
	r_1s_3	tata	r_2s_3	tada	r_3s_3	aa



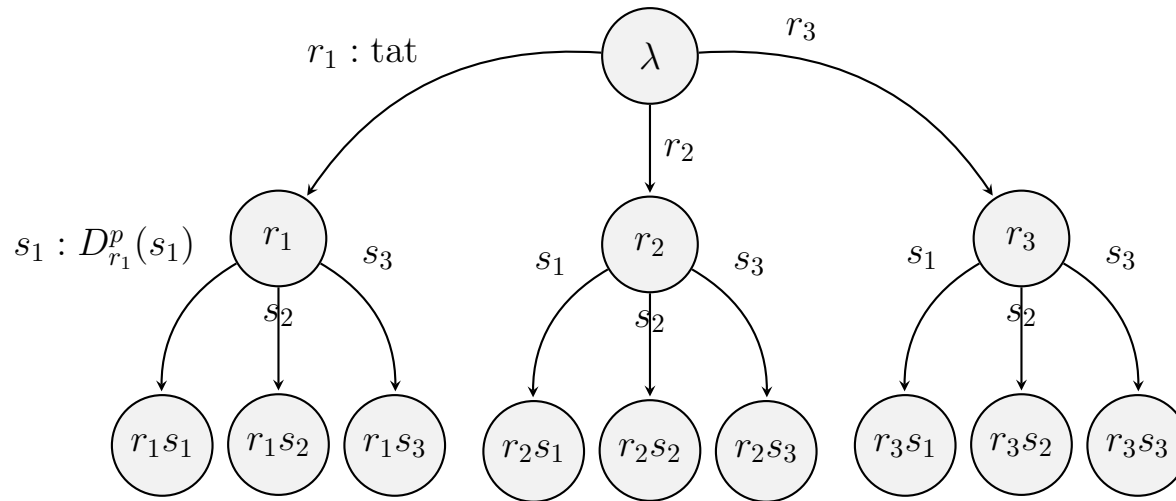
The procedure

$m : D_w^p(m)$	$r_1 s_1$ <u>tatta</u>	$r_2 s_1$ tadda	$r_3 s_1$ ata
	$r_1 s_2$ <u>tatda</u>	$r_2 s_2$ tadda	$r_3 s_2$ ada
	$r_1 s_3$ <u>tata</u>	$r_2 s_3$ tada	$r_3 s_3$ aa



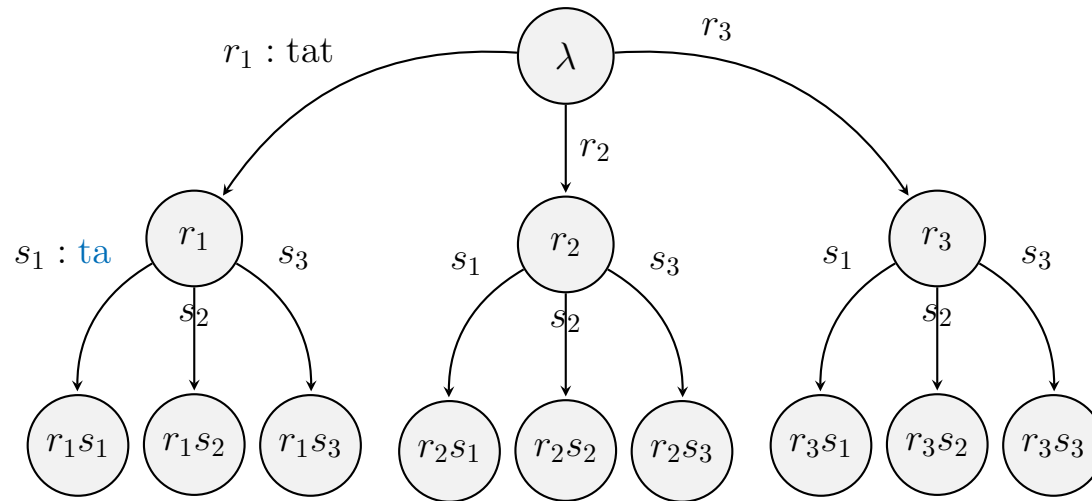
The procedure

	r_1s_1 tatta	r_2s_1 tadda	r_3s_1 ata
$m : D_w^p(m)$	r_1s_2 tatda	r_2s_2 tadda	r_3s_2 ada
	r_1s_3 tata	r_2s_3 tada	r_3s_3 aa



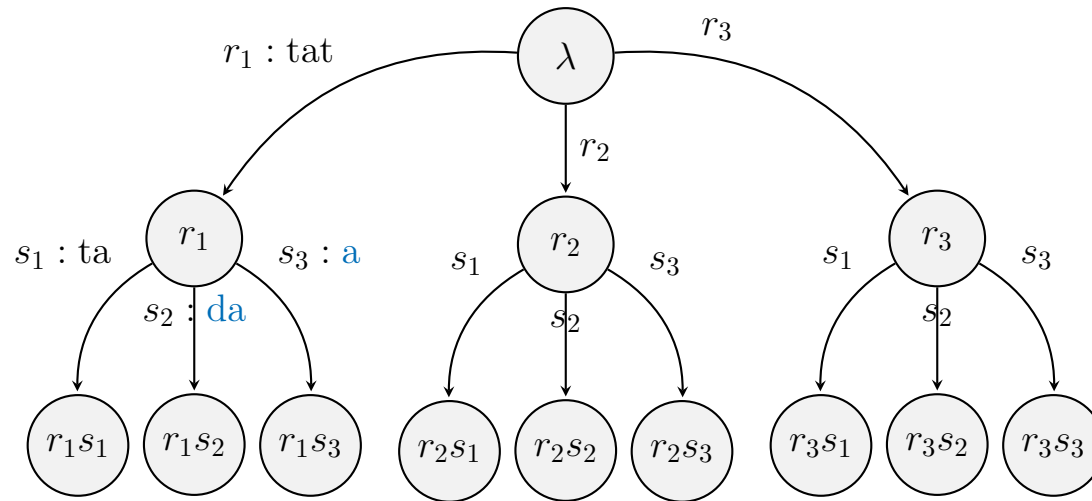
The procedure

$m : D_w^p(m)$	$r_1 s_1$ tatta	$r_2 s_1$ tadda	$r_3 s_1$ ata
	$r_1 s_2$ tatda	$r_2 s_2$ tadda	$r_3 s_2$ ada
	$r_1 s_3$ tata	$r_2 s_3$ tada	$r_3 s_3$ aa



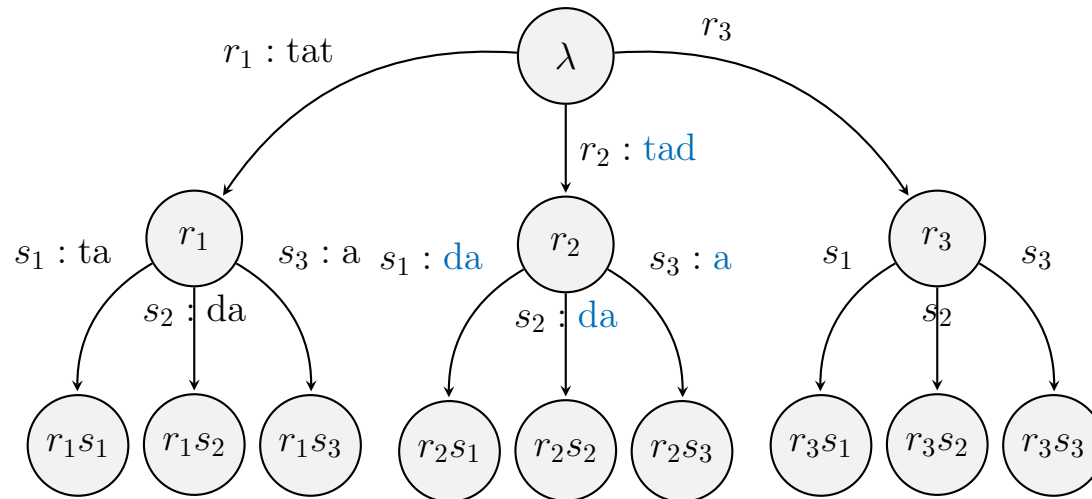
The procedure

$m : D_w^p(m)$	$r_1 s_1$ tatta	$r_2 s_1$ tadda	$r_3 s_1$ ata
	$r_1 s_2$ tat <u>da</u>	$r_2 s_2$ tadda	$r_3 s_2$ ada
	$r_1 s_3$ tata <u>a</u>	$r_2 s_3$ tada	$r_3 s_3$ aa



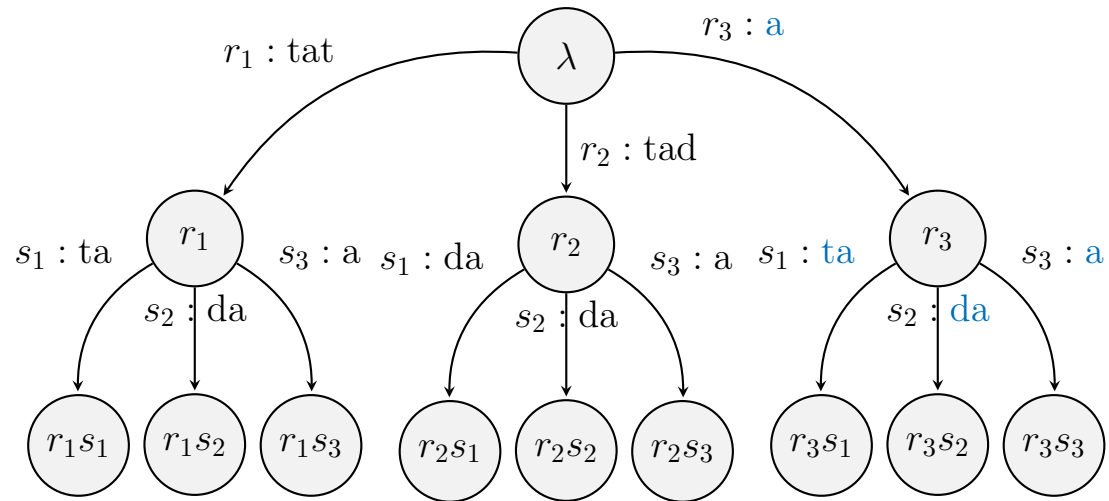
The procedure

$m : D_w^p(m)$	$r_1 s_1$	tatta	$r_2 s_1$	<u>tad da</u>	$r_3 s_1$	ata
	$r_1 s_2$	tatda	$r_2 s_2$	<u>tad da</u>	$r_3 s_2$	ada
	$r_1 s_3$	tata	$r_2 s_3$	<u>tad a</u>	$r_3 s_3$	aa



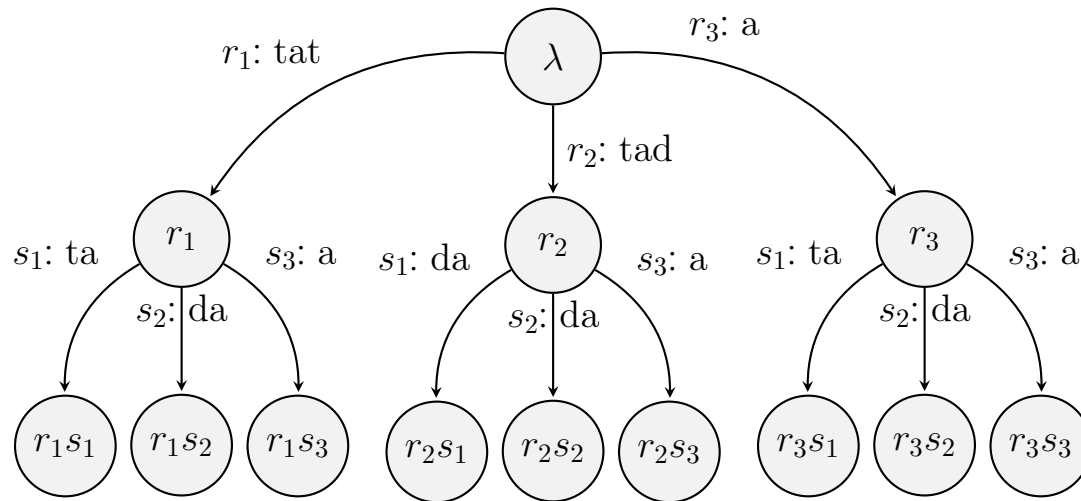
The procedure

$m : D_w^p(m)$	r_1s_1	tatta	r_2s_1	tadda	r_3s_1	<u>ata</u>
	r_1s_2	tatda	r_2s_2	tadda	r_3s_2	<u>ada</u>
	r_1s_3	tata	r_2s_3	tada	r_3s_3	<u>aa</u>



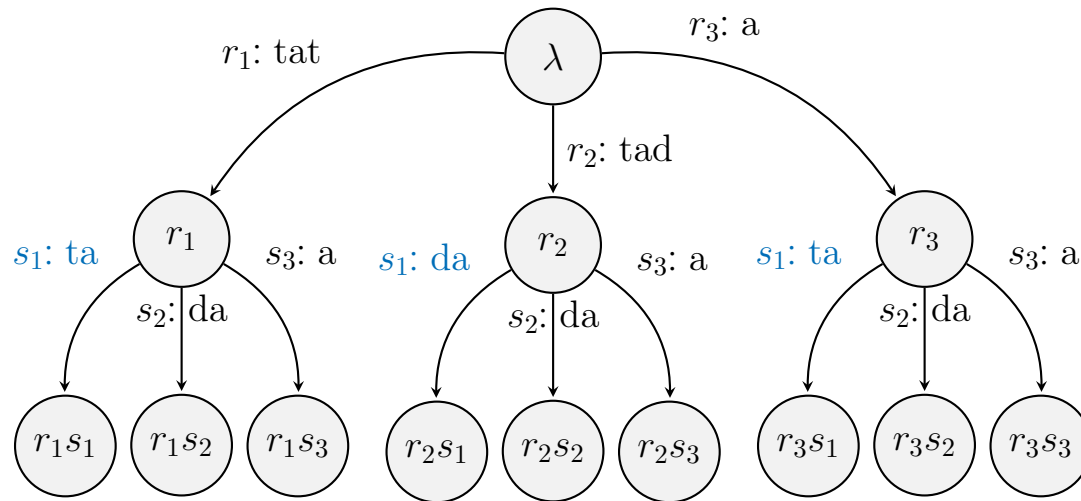
The procedure

r_1s_1	tatta	r_2s_1	tadda	r_3s_1	ata
r_1s_2	tatda	r_2s_2	tadda	r_3s_2	ada
r_1s_3	tata	r_2s_3	tada	r_3s_3	aa

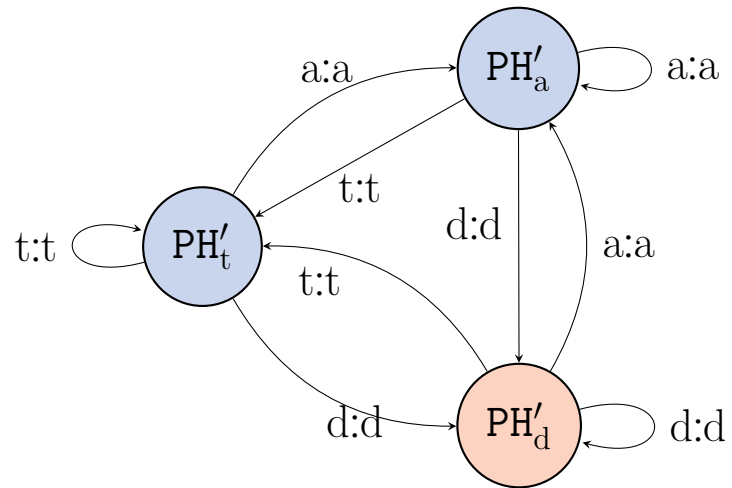


The procedure

r_1s_1	tatta	r_2s_1	tadda	r_3s_1	ata
r_1s_2	tatda	r_2s_2	tadda	r_3s_2	ada
r_1s_3	tata	r_2s_3	tada	r_3s_3	aa

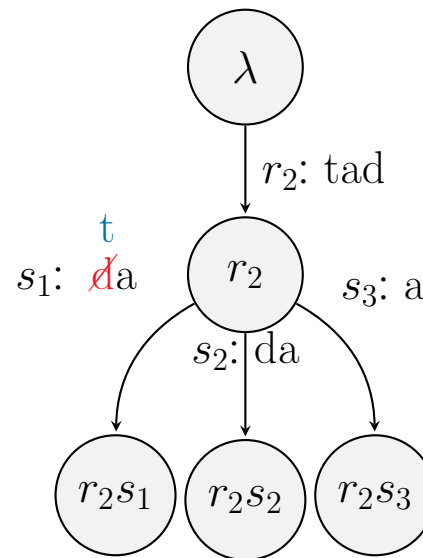
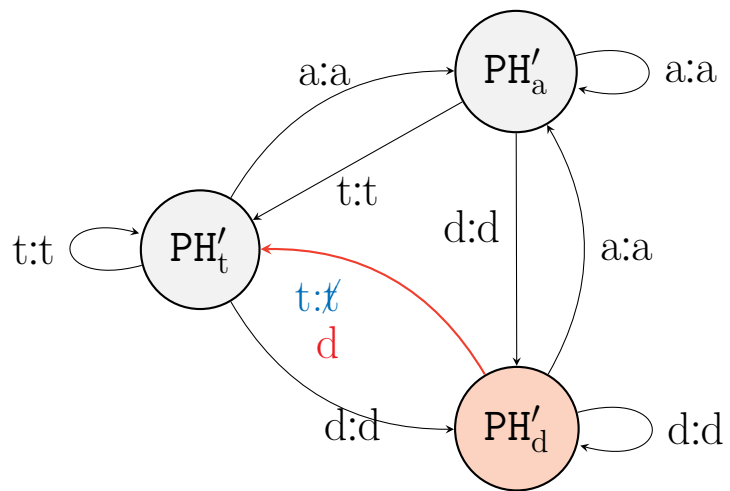


The procedure

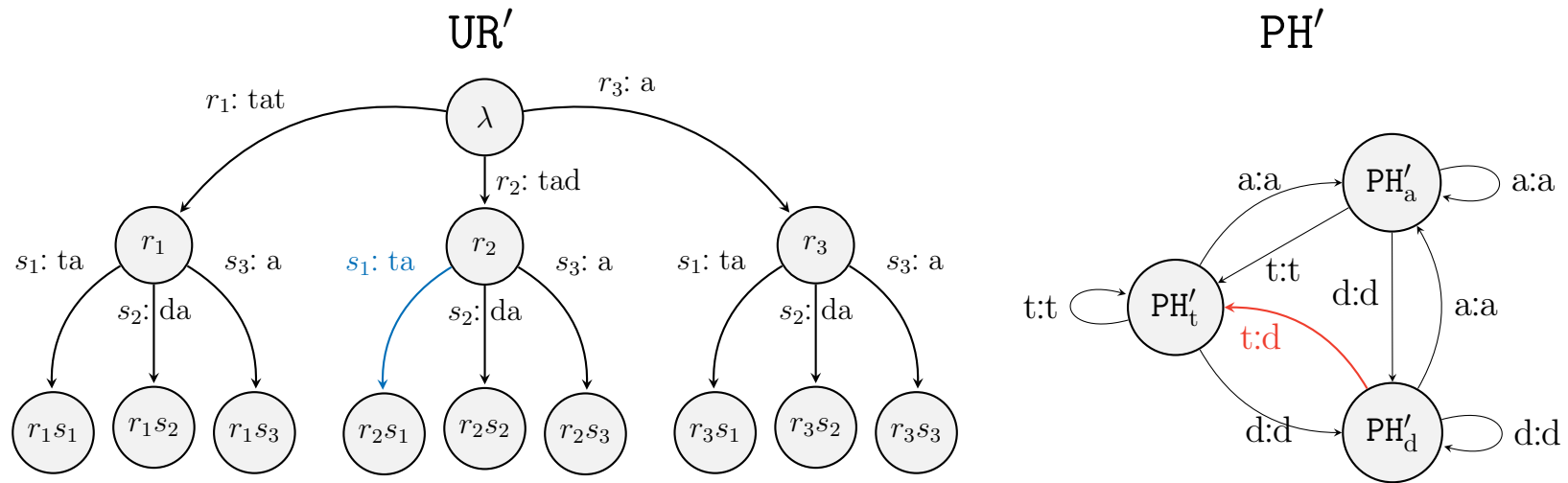


ws_1	env.	s_1
r_1s_1	tat	ta
r_3s_1	a	ta
r_2s_1	tad	da

The procedure



The procedure



$$UR'(r_1s_1) = tatta$$

$$PH'(tatta) = tatta$$

$$UR'(r_2s_1) = tadta$$

$$PH'(tadta) = tadda$$

The procedure – summary

UR	PH	Sample of PH ◦ UR
$r_1 \mapsto \text{tat} \quad s_1 \mapsto \text{ta}$		$r_1 s_1$ tatta $r_2 s_1$ tadda $r_3 s_1$ ata
$r_2 \mapsto \text{tad} \quad s_2 \mapsto \text{da}$	$t \rightarrow d / d _$	$r_1 s_2$ tatda $r_2 s_2$ tadda $r_3 s_2$ ada
$r_3 \mapsto \text{a} \quad s_3 \mapsto \text{a}$		$r_1 s_3$ tata $r_2 s_3$ tada $r_3 s_3$ aa

- Correct UR' and PH' from a sample of PH ◦ UR
- This dependent on the
 - the **subsequentiality** of UR and PH
 - that UR maps one morpheme to one UR
 - that PH is **ISL₂**

The procedure – summary

UR	PH	Sample of PH ◦ UR
$r_1 \mapsto \text{tat}$ $s_1 \mapsto \text{ta}$		r_1s_1 tatta r_2s_1 tadda r_3s_1 ata
$r_2 \mapsto \text{tad}$ $s_2 \mapsto \text{da}$	$t \rightarrow d / d _$	r_1s_2 tatda r_2s_2 tadda r_3s_2 ada
$r_3 \mapsto \text{a}$ $s_3 \mapsto \text{a}$		r_1s_3 tata r_2s_3 tada r_3s_3 aa

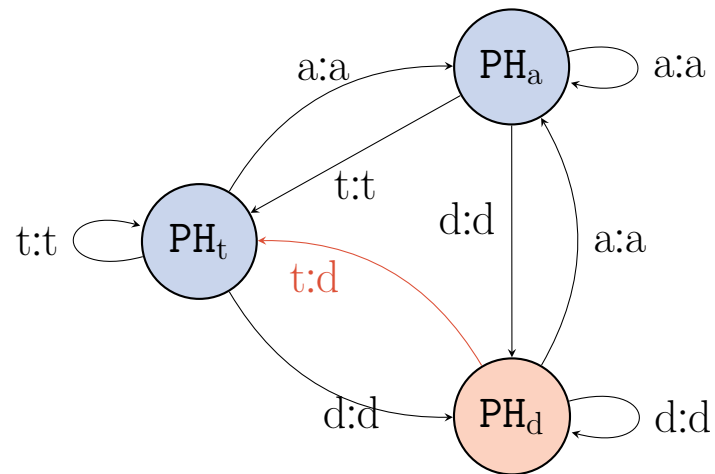
- A learner based on this procedure can learn ISL₂:
 - progressive assim./dissim.
 - regressive assim./dissim.
 - deletion
 - epenthesis
- This includes opacity (self-counterbleeding)
- So far we are limited to **single processes**

The procedure – summary

UR	PH	Sample of $PH \circ UR$
$r_1 \mapsto \text{tat} \quad s_1 \mapsto \text{ta}$		r_1s_1 tatta r_2s_1 tadda r_3s_1 ata
$r_2 \mapsto \text{tad} \quad s_2 \mapsto \text{da}$	$t \rightarrow d / d _$	r_1s_2 tatda r_2s_2 tadda r_3s_2 ada
$r_3 \mapsto \text{a} \quad s_3 \mapsto \text{a}$		r_1s_3 tata r_2s_3 tada r_3s_3 aa

- It requires that the UR is **recoverable** from $PH \circ UR$
- What are the constraints on PH? On $PH \circ UR$?

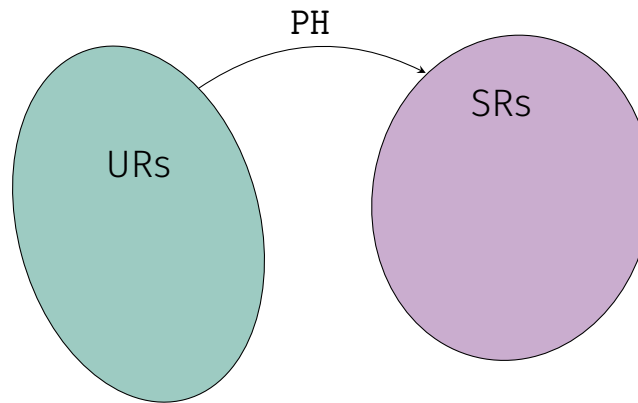
The procedure – summary



- Environment functions PH_w must split into **change** and **elsewhere** functions
- Any change PH makes must be seen at morpheme boundaries
- Formalizing these constraints on $\text{PH} \circ \text{UR}$ is work in progress

Discussion

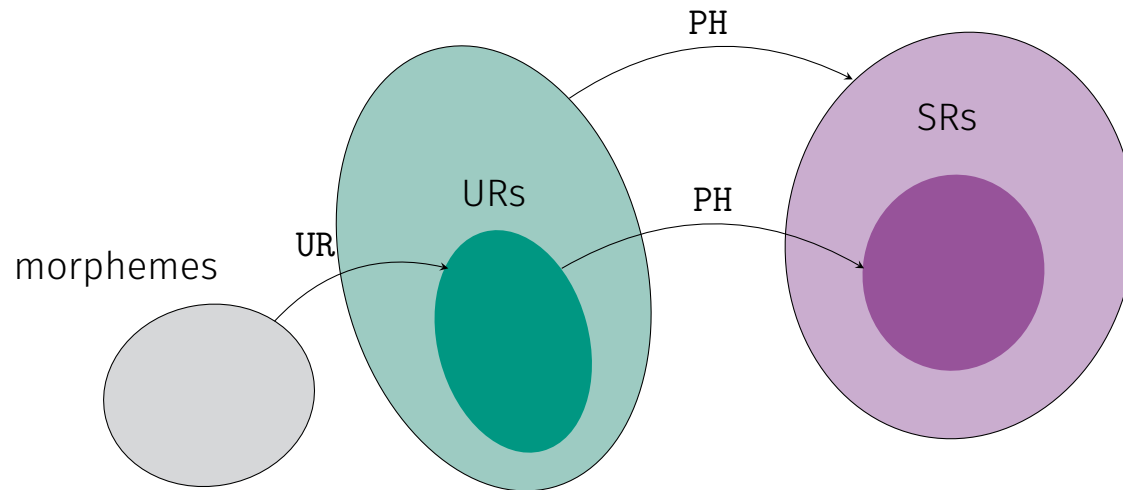
What is the nature of phonology?



What is the nature of:

- maps from URs to SRs?
- relation between SRs and URs?

What is the nature of phonology?



Assuming

- **subsequential** maps from URs to SRs, and
 - a (relatively) **concrete** relation between SRs and URs
- ...allows for a procedure for learning URs and a grammar

Future work

- Formalizing “relatively concrete”
- Extending to cases in which PH is...
 - ISL_2
 - ISL_k for some k
 - in any subsequential class with a shared structure (Jardine et al., 2014)
 - **output**-strictly local (Chandlee et al., 2015)
- Extending to...
 - featural learning (Heinz and Koirala, 2010; Chandlee et al., 2019)
 - optional/gradient processes (Shibata and Heinz, 2019; Beros and de la Higuera, 2016)

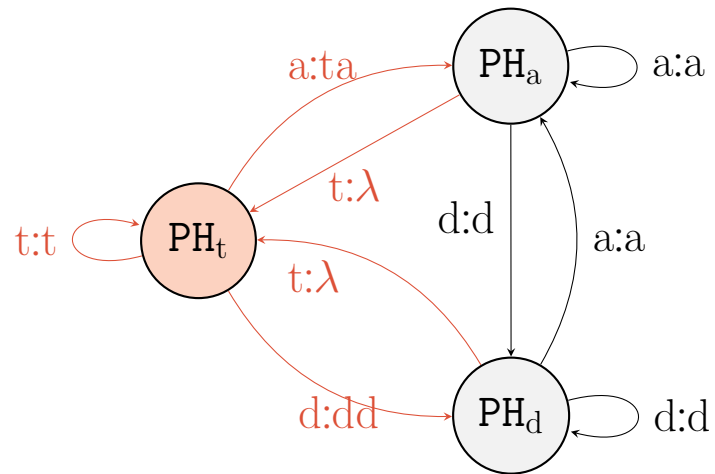
Acknowledgements

Thank you for having me!

...and many thanks to **Wenyue Hua** and **Huteng Dai**, attendees of the Rutgers/SBU/Haverford/Delaware subregular phonology workshop, **the Rutgers MathLing group**, an audience at NECPhon, and in particular **Jeff Heinz, Charles Reiss, Bruce Tesar, Adam McCollum**, and **Colin Wilson** for their insightful comments.

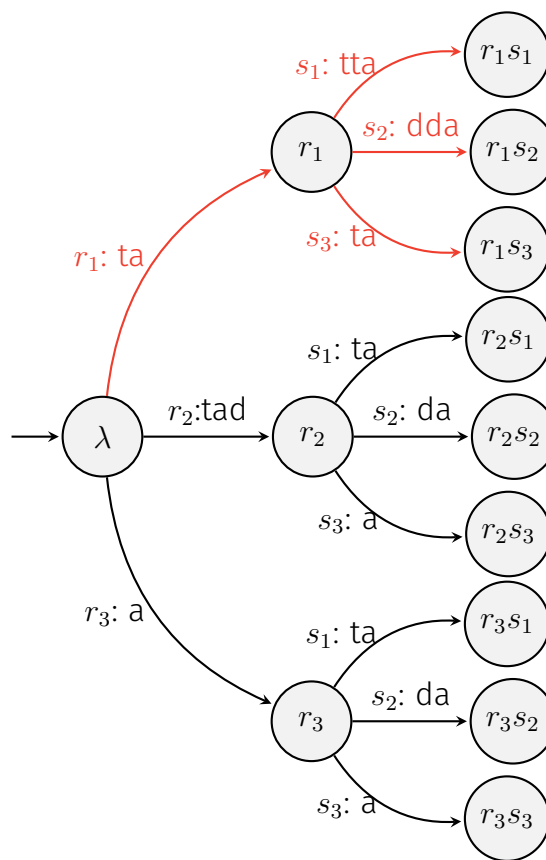
Appendix: Regressive assimilation

w	$UR(w)$	$PH \circ UR(w)$
r_1s_1	tatta	tatta
r_1s_2	tatda	tadda
r_1s_3	tata	tata
r_2s_1	tadda	tadda
r_2s_2	tadta	tadta
r_2s_3	tada	tada
r_3s_1	ata	ata
r_3s_2	ada	ada
r_3s_3	aa	aa

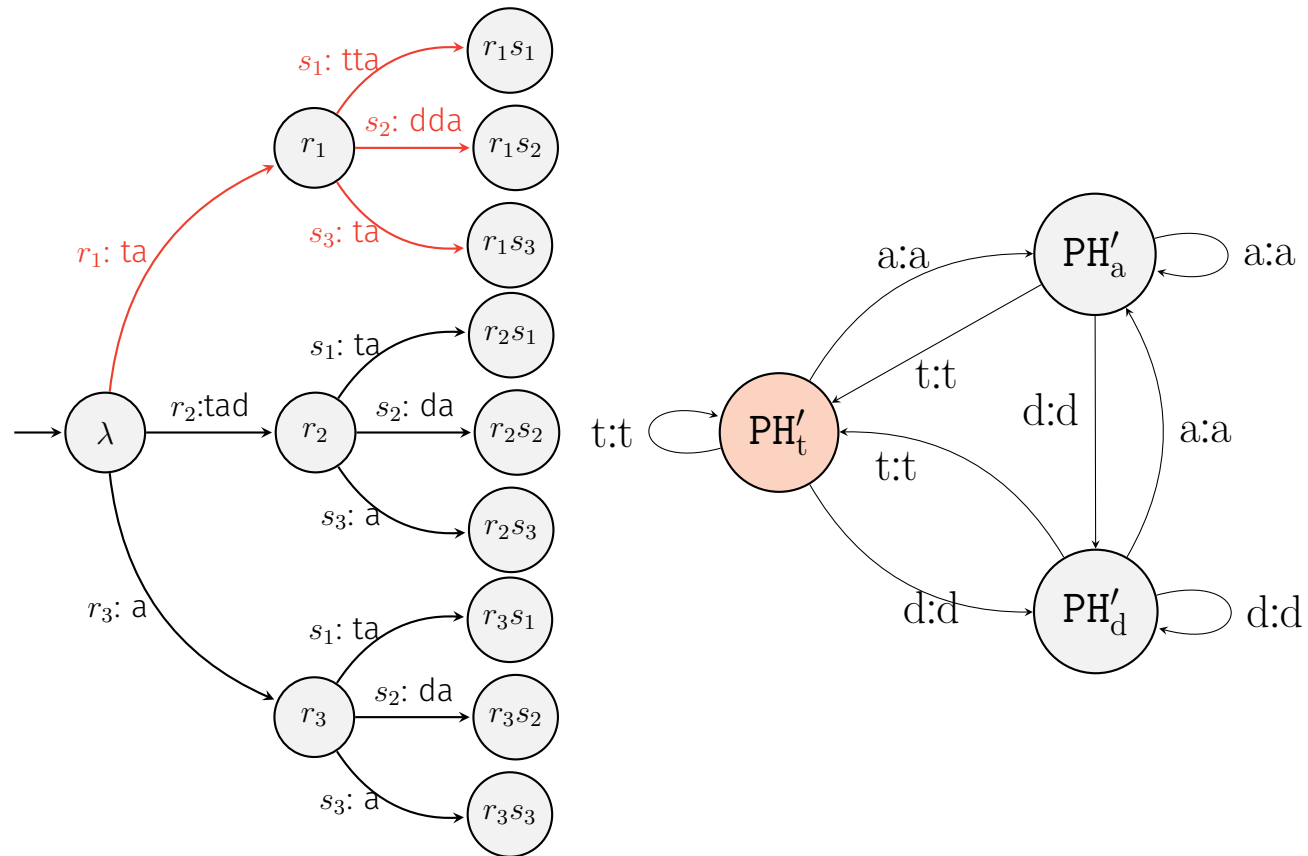


Appendix: Regressive assimilation

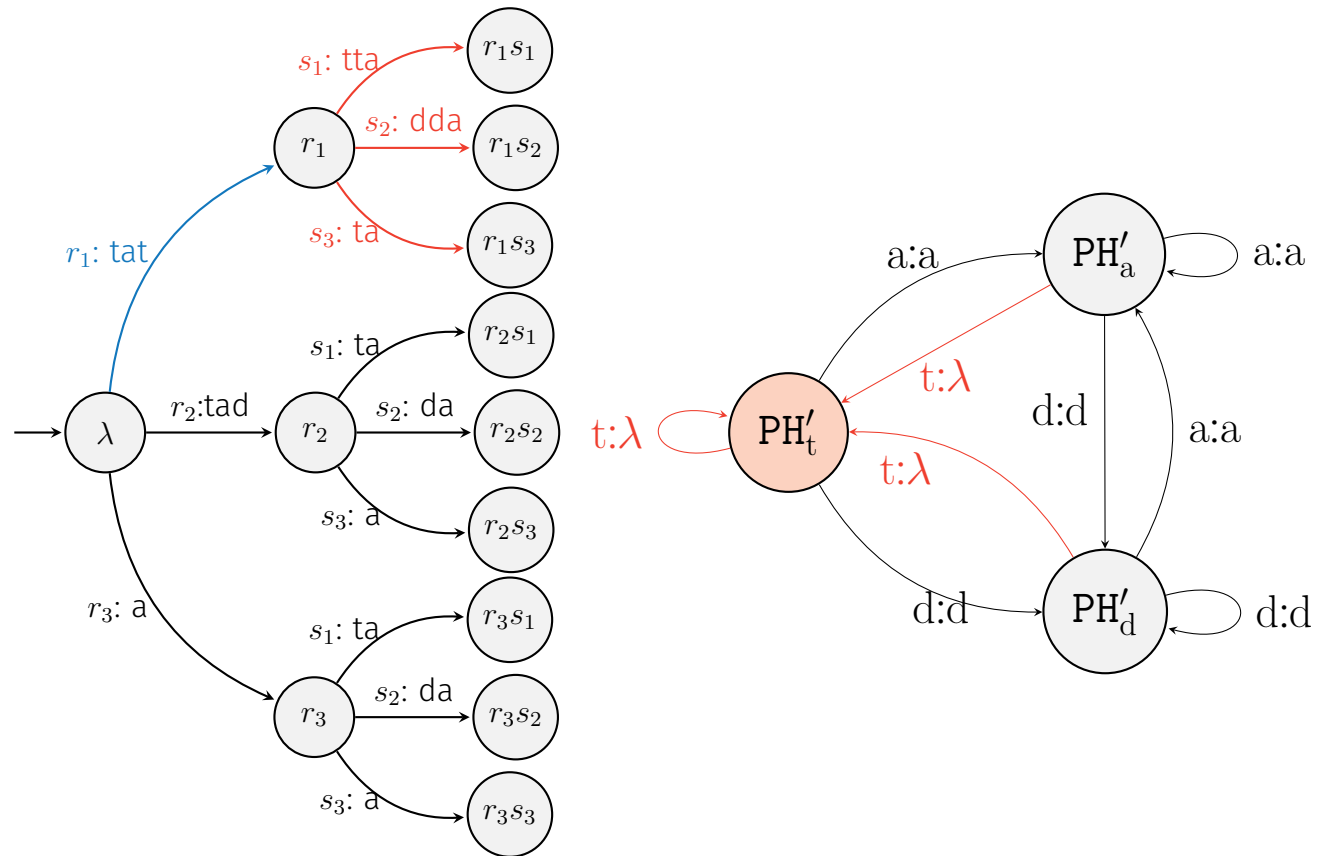
w	UR(w)	PH \circ UR(w)
r_1s_1	tatta	tatta
r_1s_2	tatda	tadda
r_1s_3	tata	tata
r_2s_1	tadda	tadda
r_2s_2	tadta	tadta
r_2s_3	tada	tada
r_3s_1	ata	ata
r_3s_2	ada	ada
r_3s_3	aa	aa



Appendix: Regressive assimilation



Appendix: Regressive assimilation



Appendix: Regressive assimilation

