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A tour of monadic dividing lines

Samuel Braunfeld (Charles University) Joint work with Chris Laskowski

September 20, 2022 UW-Madison Logic Seminar

OVERVIEW



Introduction



2 Monadic NIP



3 More monadic properties



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DEFINITION AND EXAMPLES

Definition

A theory *T* is *monadically* \mathcal{P} if any expansion of *T* by arbitrarily many unary predicates remains \mathcal{P} .

• We will consider the following hierarchy: monadically NIP, monadically stable, monadically NFCP, cellular.

Examples

- Cellular: An equivalence relation consisting of infinitely many finite classes.
- Monadically NFCP: Bounded-degree graphs.
- Monadically stable: Refining equivalence relations.
- $\bullet\,$ Monadically NIP: $(\mathbb{Q},<)$ and various tree-like structures.
- Not monadically NIP: Essentially anything with a non-unary function (e.g. vector spaces), cross-cutting equivalence relations, and the generic permutation.

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MORE ON MONADIC PROPERTIES

- Nothing beyond monadic NIP.
- Give significant information beyond their non-monadic counterparts.
- Several applications in the combinatorics of hereditary classes and countable structures.
- Preserved by taking substructure, and so about universal theories.
- Actually agree with non-monadic counterparts in universal theories [BL22b].
- We will focus on three characterizations: independence relations, decompositions, and canonical obstructions.

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CHARACTERIZATIONS

- Applications: Sparse graph classes [AA14], twin-width for hereditary (ordered) graph classes [ST21]
- We characterize monadically NIP theories in the following ways [BL21], building on [BS85] and [She86].
- The behavior of independence
- A forbidden configuration
- O Decompositions of models
- Type counting/width
- The behavior of indiscernibles
 - These are all characterizations of the theory *T* itself rather than unary expansions.
 - An intuition is that models of monadically NIP theories are 1-dimensional, or alternatively are order-like (or tree-like).

INDEPENDENCE: THE F.S. DICHOTOMY

• Finite satisfiability gives a (possibly asymmetric) notion of independence in any theory.

Definition ([She86])

Let $A
ightharpoonup_{M}^{fs} B$ mean that tp(A/MB) is finitely satisfiable in M. A theory T has the f.s.-dichotomy if given $A
ightharpoonup_{M}^{fs} B$, then for any c, $cA
ightharpoonup_{M}^{fs} B$ or $A
ightharpoonup_{M}^{fs} Bc$.

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FORBIDDEN CONFIGURATION: AN INFINITE GRID

Lemma (mostly [She86])

If T does not have the f.s. dichotomy, then some model of T codes an infinite grid (on tuples).



- So monadically NIP \Rightarrow f.s. dichotomy.
- If *T* is a universal theory then we can code grids by a boolean combination of existentials [BL22b].

DECOMPOSABILITY: *M*-F.S. SEQUENCES

Definition

Given a model M, $(A_i : i \in I)$ is an M-f.s. sequence if $A_i \bigcup_{M}^{f_s} \{A_{\leq i}\}$.

A *linear decomposition* of *N* is a partition $N = \bigsqcup_i A_i$ and a model *M* (not necessarily in *N*) such that $(A_i : i \in I)$ is an *M*-f.s. sequence.

Lemma ([She86])

If T has the f.s. dichotomy, then any partial linear decomposition of $N \models T$ can be extended to a full linear decomposition of N.

• The f.s. dichotomy is exactly what we need to inductively extend, one point at a time.

MONADIC STABILITY

Theorem ([BS85])

For a complete theory T, the following are equivalent.

- **1** *T* is monadically stable
- **2** *T* is monadically NIP and stable
- Models of T are tree-decomposable
- *T* is stable and forking is transitive on singletons and totally *trivial*.
 - Totally trivial forking: If $A \not\perp_C B$, then there are $a \in A, b \in B$ such that $a \not\perp_C b$. (Vector spaces are a non-example.)

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MORE MONADIC STABILITY

- Triviality on transitivity mean forking-dependence yields on equivalence relation on singletons.
- Structure (for uncountable models): Tree decompositions
- Non-structure: the order property (by an atomic formula [NMP⁺21])
- Applications: ω-categorical classification [Lac92], sub-exponential growth rates of ω-categorical structures [Bra22]

MONADIC NFCP

- Equivalence relation with arbitrarily large finite classes as a paradigm of stable but not NFCP.
- So an equivalence relation with infinitely many infinite classes is not monadically NFCP.

Theorem (mostly [Las13])

The following are equivalent for a complete theory T.

- **1** *T is monadically NFCP.*
- T mutually algebraic.
- T is weakly minimal and forking is totally trivial.
- Models of T can be partitioned into a sunflower with countable pieces (MA-connected components).
- See next slide for non-structure characterizations.

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More on monadic NFCP

- Weak minimality: forking dependence just comes from algebraic closure.
- Mutual algebraicity: definable sets come from bounded-degree relations. In particular, after naming |L|-many (or so) constants, models are quantifier-free interdefinable with bounded-degree structures.
- Structure: bounded-degree structures naturally split into countable connected components.
- For non-structure, assume a finite language. (*T* is monadically NFCP iff every finite reduct is.)
- If a universal *T* is monadically stable but not monadically NFCP, then some models defines an infinite equivalence relation (by a universal formula) [BL22a] (building on [LT20]).
- Robust closure under expansions

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CELLULARITY

- Partitions like monadic NFCP, but with finiteness conditions.
- Does not encode a linear order or infinite partition.
- Preserved by finite monadic expansions.
- Equivalent to ω-categorical and an independence condition: ω-stable with Morley rank 1, and trivial algebraic closure [Sch90].
- If *T* is monadically NFCP but non-cellular, then every model admits an elementary extension containing infinitely many infinite new MA-connected components [BL22c].

APPLICATIONS OF CELLULARITY

- Counting countable structures of a given age [MPW92]
- Counting substructures of a countable structure [LM96]
- Counting finite structures in a hereditary class [LT22].
- Counting structures bi-embeddable with a given countable structure [BL22a].

Conjecture (extending Pouzet-Sauer-Thomassé [LPSW21])

Given an age \mathcal{A} , let $|Mod(\mathcal{A})/\equiv|$ count the bi-embeddability classes of countable structures of age \mathcal{A} . Then $|Mod(\mathcal{A})/\equiv| \in \{1, \aleph_0, \aleph_1, 2^{\aleph_0}\}$. Furthermore, these cases correspond to cellular, (monadically) stable, (monadically) NIP, and everything else.

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