Basic Notions

Term:

Constant, variable, or compound term.

Constant:

Integer or atom.

Compound Term:

Functor, arguments

 $f(t_1,\ldots,t_n)$

Functor:

Name, arity

f/n

Goal:

Atom or compound term.

Logic Programs

Clause:

Universally quantified logical sentence

 $A \leftarrow B_1, \ldots, B_k, \ k \ge 0$

A and B_i -s are goals.

Declarative reading:

A is implied by the conjunction of the B_i -s.

Procedural reading:

To answer query A, answer the conjunctive query B_1, \ldots, B_k .

Logic Program:

Finite set of clauses

Computation

Query:

Existentially quantified conjunction

 $\leftarrow A_1, \ldots, A_n, \ n > 0$

 A_i -s are goals.

Computation of a Logic Program P:

finds an instance of a given query logically deducible from P.

How to Compute

Start from initial query G. Computation terminates – success or failure. Computation does not terminate – no result. Output of a successful computation: the instance of G proved.

A given query can have several successful computations with different output.

Abstract Interpreter

INPUT:

A logic program P and a query G

OUTPUT:

 $G\theta$, if this was the instance of G deduced from P, or *failure* if failure has occurred.

ALGORITHM:

Let resolvent be G

While resolvent is not empty do

 $\underline{\text{Choose}}$ a goal A from the resolvent

<u>Choose</u> a renamed clause $A' \leftarrow B_1, \ldots, B_n$ from P such that A and A' unify with mgu θ (exit if no such goal and clause exist)

Remove A from and <u>add</u> B_1, \ldots, B_n to the resolvent

Apply θ to the resolvent and to G

If the resolvent it empty, return G, else return *failure*.

Choosing and Adding

Left unspecified in the abstract interpreter

Must be resolved in any realization of the computational model

<u>Two Choices</u>

Completely different nature

Choice of a goal:

Arbitrary

Does not affect computation If there exists a successful computation by choosing one goal, then there is a successful computation by choosing any other goal.

Choice of a clause:

Non-deterministic

Affects computation

Choosing one clause might lead to a successful computation, while choosing some other might lead to failure

Adding Goal to Resolvent Assume always the leftmost goal to be chosen Then Adding new goal to the beginning of the resolvent gives depth-first search

Adding new goal to the end of the resolvent gives breath-first search

Prolog's Solution

Choice of a goal:

leftmost

Choice of a clause:

Topmost

Adding new goal to the resolvent:

At the beginning